Information Element

octets		
1	1	length
Element ID	Length	Information

Information Element	Element ID
SSID	0
Supported rates	1
FH Parameter Set	2
DS Parameter Set	3
CF Parameter Set	4
ТІМ	5
IBSS Parameter Set	6
Country	7
Legacy Indication (11g)	8
Reserved	9-15
Challenge Text	16
Reserved for challenge text extension	17-31
Reserved	32-255

Elements



SSID Elements



- indicates the identity of an ESS or IBSS
- a '0' length information field indicates the broadcast SSID

Supported Rate Elements



- The number of supported rates is **14** (a/b/g).
- Each supported rate belonging to the BSSBasicRateSet is encoded as an octet with the msb (bit 7) set to 1
 - e.g., a 1 Mbit/s rate is encoded as X'82'
- Rates not belonging to the BSSBasicRateSet are encoded with the msb set to 0
 - e.g., a 2 Mbit/s rate is encoded as X'04'.

ERP Information Elements



Bit b0	NonERP_Present
0	No NonERP stations are within the BSS
1	There are NonERP stations within the BSS

Bit b1	Use_Protection
0	STAs with an ERP should not use protection mechanisms for MPDUs transmitted at one of the ERP-OFDM rates.
1	STAs with an ERP shall use protection mechanisms for MPDUs transmitted at one of the ERP-OFDM rates.

- transmitted from AP in BSS or STA in IBSS
- defined in IEEE 802.11g
- Protection mechanism

Use CTS frame to update the NAV of all receiving STAs prior to <u>the</u> <u>transmission of a frame that may or may not be understood by receivers.</u> The updated NAV period shall be longer than or equal to the total time required to send the data and any required response frames.

MAC Architecture



MAC Architecture

- Distributed Coordination Function (DCF)
 - The fundamental access method for the 802.11 MAC, known as Carrier Sense Multiple Access with Collision Avoidance (CSMA/CA).
 - Shall be implemented in all stations and APs.
 - Used within both ad hoc and infrastructure configurations.
- Point Coordination Function (PCF)
 - An alternative access method
 - Shall be implemented on top of the DCF
 - A point coordinator (polling master) is used to determine which station currently has the right to transmit.
 - Shall be built up from the DCF through the use of an access priority mechanism.
 - Different accesses of traffic can be defined through the use of different values of IFS.
 - Shall use a Point IFS (PIFS) < Distributed IFS (DIFS)</p>

MAC Architecture

- Point coordinated traffic shall have higher priority to access the medium, which may be used to provide a contention-free access method.
- The priority access of the PIFS allows the point coordinator to seize control of the medium away from the other stations.
- Coexistence of DCF and PCF
 - Both the DCF and PCF shall coexist without interference.
 - They are integrated in a superframe in which a contention-free burst occurs at the beginning, followed by a contention period.



Distributed Coordination Function

- Allows for automatic medium sharing between similar and dissimilar PHYs through the use of CSMA/CA and a random backoff time following a busy medium condition.
- All directed traffic uses immediate positive ack (ACK frame) where retransmission is scheduled by the sender if no ACK is received.
- The virtual Carrier Sense mechanism is achieved by distributing medium busy reservation information through an <u>exchange of</u> <u>special small RTS and CTS frames</u> (contain a duration field) prior to the actual data frame.
 - Unicast only, not used in multicast/broadcast.
- The use of RTS/CTS is under control of RTS_Threshold (payload length, under which without any RTS/CTS prefix).
- All stations are required to be able to receive any frame transmitted on a given set of rates, and must be able to transmit at (at least) one of these rates.
 - Virtual Carrier Sense mechanism still works on multiple rates environments.

Distributed Coordination Function

- Carrier Sense shall be performed both through *physical* (listen) and *virtual* mechanisms.
- Physical Carrier Sense Mechanism
 - A physical carrier sense mechanism shall be provided by the PHY.
- Virtual Carrier Sense Mechanism
 - Provided by the MAC, named Net Allocation Vector (NAV), which maintains a prediction of future traffic based on duration information announced in RTS/CTS frames.
- MAC-Level Acknowledgments (Positive Acknowledgment)
 - To allow detection of a lost or errored frame, an ACK frame shall be returned immediately following a successfully received frame.
 - The gap between the received frame and ACK frame shall be SIFS.
 - The frame types should be acknowledged with an ACK frame:
 - » Data
 - » Poll

Why does not Ethernet use ACK frame?

» Request» Response

The lack of an ACK frame means that an error has occurred.

Distributed Coordination Function --Inter-Frame Space (IFS)

- A station shall determine that the medium is free through the use of carrier sense function for the interval specified.
- Three different IFS's are defined to provide priority levels.
- Short-IFS (SIFS)
 - Shall be used for an ACK frame, a CTS frame, by a station responding to any polling.
 - It may also be used by a PC for any types of frames during the CFP.
 - Any STA intending to send only these frame types shall be allowed to transmit after the SIFS time has elapsed following a busy medium.
- PCF-IFS (PIFS)
 - Shall be used only by the PCF to send any of the Contention Free Period frames.
 - The PCF shall be allowed to transmit after it detects the medium free for the period PIFS, at the start of and during a CF-Burst.
- DCF-IFS (DIFS)
 - Shall be used by the DCF to transmit asynchronous MPDUs.
 - A STA using the DCF is allowed to transmit after it detects the medium free for the period DIFS, as long as it is not in a backoff period.
- Extended IFS (EIFS), (1ms)

Time Intervals SIFS/PIFS/DIFS



 $\begin{array}{l} D1 = aRxRFDelay + aRxPLCPDelay (referenced from the end of the last symbol of a frame on the medium) \\ D2 = D1 + Air Propagation time \\ Rx/Tx = aRXTXTurnaroundTime (begins with a PHYTXSTART.request) \\ M1 = M2 = aMACPrcDelay \\ CCAdel = aCCATime - D1 \end{array}$

EIFS

- The EIFS shall begin following indication by the PHY that the medium is idle after detection of the erroneous frame, without regard to the virtual carrier-sense mechanism.
- The EIFS is defined to provide enough time for another STA to acknowledge what was, to this STA, an incorrect received frame before this STA commences transmission.
- EIFS = aSIFSTime + (8×ACKsize) + aPreambleLength + PLCPHeaderLength + DIFS,

where ACKsize is computed based on 1Mbps data rate.

Distributed Coordination Function --Random Backoff Time

- Before transmitting asynchronous MPDUs, a STA shall use the carrier sense function to determine the medium state.
- If busy, the STA shall defer until after a DIFS gap is detected, and then generate a random backoff period for an additional deferral time (resolve contention).

Backoff time = Random() * Slot time

Random() = Pseudorandom integer drawn from a uniform distribution over the interval [0, CW]. CW = An integer between CWmin and CWmax Slot Time = Transmitter turn-on delay + medium propagation delay + medium busy detect response time

Binary Exponentional Backoff Window



15~1023 for FHSS PHY

Source: IEEE Std 802.11-1997

14.8.2 FH PHY attributes: Table 49

63~1023 for IR PHY

Source: IEEE Std 802.11-1997

16.4 PHY attributes: Table 74

31~1023 for DSSS PHY Source: IEEE Std 802.11-1997 15.3.2 DSSS PHY MIB: Table 58

15~1023 for DSSS ERP PHY (>20Mb/s) 31 ~1023 for DSSS ERP PHY (≤20Mb/s) Source: IEEE Std 802.11g-2001 19.4.3.8.5 PHY Page 12

Basic Access Protocol Features

- Use Distributed Coordination Function (DCF) for efficient medium sharing without overlap restrictions.
 - Use CSMA with Collision Avoidance derivative.
 - Based on Carrier Sense function in PHY called Clear Channel Assessment (CCA).
- Robust for interference (use positive acknowledge).
 - CSMA/CA + ACK for unicast frames, with MAC level recovery.
 - CSMA/CA for Broadcast frames.
- Parameterized use of RTS / CTS to provide a *Virtual Carrier Sense* function to protect against *Hidden Nodes*.
 - Duration information is distributed by both transmitter and receiver through separate RTS and CTS Control Frames.
- Includes fragmentation to cope with different PHY characteristics.
- Frame formats to support the access scheme
 - Infrastructure and Ad-Hoc Network support.
 - Wireless Distribution System.

CSMA/CA Explained



- **Reduce collision probability** where mostly needed.
 - Stations are waiting for medium to become free.
 - Select Random Backoff after a Defer, resolving contention to avoid collisions.
- Efficient Backoff algorithm stable at high loads.
 - Exponential Backoff window increases for retransmissions.
 - Backoff timer elapses only when medium is idle.
- Implement different fixed priority levels.
 - To allow immediate responses and PCF coexistence.

CSMA/CA + ACK protocol



Defer access based on Carrier Sense.

- CCA from PHY and Virtual Carrier Sense state.

- Direct access when medium is sensed free longer than DIFS, otherwise defer and backoff.
- Receiver of directed frames to return an ACK immediately when CRC correct.
 - When no ACK received then retransmit frame after a random backoff (up to maximum limit).

Throughput Efficiency



- Efficient and stable throughput.
 - Stable throughput at overload conditions.
 - To support Bursty Traffic characteristics.

Hidden Node Problem

• Transmitters contending for the medium may not "Hear each other" as shown below.



 Separate Control frame exchange (RTS / CTS) between transmitter and receiver will *Reserve the Medium* for subsequent data access.

- Duration is distributed around both Tx and Rx station.

Hidden Node Provisions



- Duration field in RTS and CTS frames distribute Medium Reservation information which is stored in a Net Allocation Vector (NAV).
- Defer on either NAV or "CCA" indicating *Medium Busy*.
- Use of RTS / CTS is optional but <u>must</u> be implemented.
- Use is controlled by a *RTS_Threshold* parameter per station.
 - To limit overhead for short frames. (200 bytes)

RTS/CTS Overhead Impact



Good mixed Throughput (long inbound frames) efficiency.

Distributed Coordination Function --DCF Access Procedure

- Backoff Procedure
 - A backoff time is selected first. The Backoff Timer shall be frozen while the medium is sensed busy and shall decrement only when the medium is free (resume whenever free period > DIFS).
 - Transmission whenever the Backoff Timer reaches zero.
 - A STA that has just transmitted a frame and has another frame ready to transmit (queued), shall perform the backoff procedure (fairness concern).
 - Tends toward fair access on a FCFS basis.



Distributed Coordination Function --DCF Access Procedure

- **RTS/CTS Recovery Procedure and Retransmit Limits**
 - After an RTS is transmitted, if the CTS fails in any manner within a predetermined CTS_Timeout (T1), then a new RTS shall be generated (the CW shall be doubled).
 - This procedure shall continue until the RTS_Re-Transmit_Counter reaches an RTS_Re-Transmit_Limit.
 - The same backoff mechanism shall be used when no ACK is received within a predetermined ACK_Window(T3) after a directed DATA frame has been transmitted.
 - This procedure shall be continue until the ACK_Re-Transmit_Counter reaches an ACK_Re-Transmit_Limit.
 - STAs shall maintain a short retry count (for MAC frame <= RTS_Threshold) and a long retry count (for MAC frame > RTS_Threshold) for each MSDU and MMPDU awaiting transmission.
 - » These counts are incremented and reset independently of each other.

Distributed Coordination Function --Fragment

- Control of the Channel (Fragment)
 - The IFS is used to provide an efficient MSDU delivery mechanism.
 - Once a station has contended for the channel, it will continue to send fragments until either all fragments of a MSDU have been sent, an ack is not received, or the station can not send any additional fragments due to <u>a dwell time boundary</u>.
 - If the source station does not receive an ack frame, it will attempt to retransmit <u>the fragment</u> at a later time (according to the backoff algorithm).
 - When the time arrives to retransmit the fragment, the source station will contend for access in the contention window.



Distributed Coordination Function --DCF Access Procedure

- **RTS/CTS Usage with Fragmentation**
 - The RTS/CTS frames define the duration of the first frame and ack.
 - The duration field in the data and ack frames specifies the total duration of the next fragment and ack.
 - The last Fragment and ACK will have the duration set to zero.
 - Each Fragment and ACK acts as a virtual RTS and CTS.
 - In the case where an ack is not received by the source station, the NAV will be marked busy for next frame exchange.
 - » This is the worst case situation.
 - If the ack is not sent by the destination, stations that can only hear the destination will not update their NAV and be free to access the channel.
 - All stations that hear the source will be free to access the channel after the NAV from Fragment 1 has expired.
 - The source must wait until the NAV (Fragment 1) expires before attempting to contend for the channel after not receiving the ack.

RTS/CTS Usage with Fragmentation



Fragmentation (1/2)



- Burst of Fragments which are individually acknowledged.
 - For Unicast frames only.
- Random backoff and retransmission of failing fragment when no ACK is returned.
- Duration information in data fragments and ACK frames causes NAV to be set, for medium reservation mechanism.

Fragmentation (2/2)

- The length of a fragment MPDU shall be an equal number of octets for all fragments except the last, which may be smaller.
- The length of a fragment shall never be larger than aFragmentationThreshold unless WEP is invoked for the MPDU.
 - Because the MPDU shall be expanded by IV and ICV.
- The sequence number shall remain the same for all fragments of a MSDU or MMPDU.
- The fragments shall be sent in order of lowest fragment number to highest fragment number (start at zero, and increased by one).
- More Fragments bit is used to indicate the last (or only) fragment of the MSDU or MMPDU.

Defragmentation

- The header of each fragment contains the following information that is used by the destination STA to reassemble the MSDU or MMPDU.
 - Frame type.
 - Address of the sender.
 - Destination address.
 - Sequence Control field.
- More Fragments indicator. If WEP has been applied, it shall be decrypted before the defragmentation.
- All STAs shall maintain a Receive Timer for each MSDU or MMPDU. If the timer is not maintained, all the fragments belong to the part of an MSDU or MMPDU are discarded.
- If the receive MSDU timer exceeds aMaxReceiveLifetime, then all received fragments of this MSDU or MMPDU are discarded.

DCF -- Broadcast and multicast

- Broadcast and multicast MPDU transfer procedure
 - In the absence of a PCF, when broadcast or multicast MPDUs are transferred from a STA with the ToDS bit clear, only the basic access procedure shall be used. Regardless of the length of the frame, no RTS/CTS exchange shall be used. (will not receive ACK)
 - Any broadcast or multicast MPDUs transferred from a STA with a ToDS bit set shall obey the rules for RTS/CTS exchange, because the MPDU is directed to the AP (will receive ACK).
 - This no MAC-level recovery on broadcast or multicast frames, except for those frames sent with ToDS bit set.
 - The broadcast/multicast message shall be distributed into the BSS, so the STA originating the message will also receive the message. Therefore, all STAs must filter out broadcast/multicast messages that contain their address as the source address.
 - Broadcast/multicast MSDUs shall be propagated throughout the ESS.